



# Linear time invariant convex optimal control

$$\min_{x, u} J = \sum_{t=0}^{\infty} \ell(x(t), u(t))$$

$x$ : state  $u$ : action

$$\text{St: } u(t) \in U \quad x(t) \in X$$

$$x(t+1) = A x(t) + B u(t)$$

$$x(t) \in \mathbb{R}^n \quad u(t) \in \mathbb{R}^m$$

$$x(0) = z$$

variables:  $x(0), x(1), \dots \in \mathbb{R}^n$  can think  $x(t)$  as variable and  $x(t+1) = A x(t) + B u(t)$  as equality constraints  
 $u(0), u(1), \dots \in \mathbb{R}^m$  or think  $x(t+1) = A x(t) + B u(t)$  as recursion expression

problem data:

dynamics and input matrices  $A \in \mathbb{R}^{n \times n}$   $B \in \mathbb{R}^{n \times m}$

convex stage cost function  $\ell: \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R}$   $\ell(0,0) = 0$

convex state set and convex input set  $X, U \quad 0 \in X \quad 0 \in U$

Greedy control:  $u(t) = \operatorname{argmin} \{ \ell(x(t), w) \mid w \in U, A x(t) + B w \in X \}$   
 ignores the cost in the future, except for  $x(t+1) \in X$

Greedy control typically works poorly

But it can work very well:

if  $A$  has very small spectral radius / small norm

then  $x(t+1)$  does not depend much on  $x(t)$ , the dynamic system has "strongly fading memory"

## Solution via dynamic programming (Reinforcement learning !!!)

the optimal value of the convex optimization problem is a convex function of  $z$

(Bellman) value function  $V(z)$  is the optimal value of control problem as a function of initial state  $z$

$V$  satisfies Bellman or dynamic programming equation

$$V(z) = \inf_w \{ \underbrace{\ell(z, w)}_{\text{immediate cost}} + \underbrace{V(Az + Bw)}_{\text{minimum future cost}} \mid w \in U, Az + Bw \in X \}$$

optimal  $u$  given by

$$u^*(t) = \operatorname{argmin}_{\substack{w \in U \\ A x(t) + B w \in X}} \ell(x(t), w) + V(A x(t) + B w)$$

the optimal input is a function of state  $u^*(t) = \phi(x(t))$  State fed back form

eg. Linear quadratic regulator

special case of linear convex optimal control with

$$U \in \mathbb{R}^m \quad X \in \mathbb{R}^n$$

$$\ell(X(t), U(t)) = X(t)^T Q X(t) + U(t)^T R U(t) \quad Q \geq 0 \quad R \geq 0$$

Solved using dynamic programming

$$\text{value function } V(z) = \min_{u, x} \sum_{t=0}^{\infty} \ell(X(t), U(t)) = z^T P z$$

minimizing a quadratic function over a subset of variables

gives a quadratic function (described by scalar component)

P can be found by solving an algebraic Riccati equation (ARE)

$$P = Q + A^T P A - A^T P B (R + B^T P B)^{-1} B^T P A$$

## • Finite horizon approximation

Use finite horizon  $T$ , impose terminal constraint  $X(T) = 0$

$$\min_{u, x} \sum_{t=0}^{T-1} \ell(X(t), U(t))$$

$$U(t) \in U \quad X(t) \in X$$

$$X(t+1) = A X(t) + B U(t)$$

$$X(0) = z \quad X(T) = 0$$

gives a suboptimal input of the original optimal input problem

## • Model predictive control (MPC)

at each time  $t$ , solve the (planning) problem

$$\min_{u, x} \sum_{\tau=t}^{t+T-1} \ell(X(\tau), U(\tau))$$

$$\text{s.t. } U(\tau) \in U \quad X(\tau) \in X$$

$$X(\tau+1) = A X(\tau) + B U(\tau)$$

$$X(t+T) = 0$$

do a complete planning exercise, use the first input, and re-plan

this gives a complicated state feedback control  $U(t) = \phi_{\text{MPC}}(X(t))$

## • Variations on MPC

— add final state cost  $\hat{v}(X(T))$  instead of insisting on  $X(t+T) = 0$   
if  $\hat{v} = v$ , MPC gives optimal input

— convert hard constraints to violation penalties  
avoids problem of planning problem infeasibility

— solve MPC every  $k$  steps ( $k \geq 1$ )

## • Explicit MPC

MPC with  $l$  quadratic,  $X$  and  $U$  polyhedral  
can show change is piece-wise affine:

$$\Delta u_{\text{mpc}}(z) = K_j z + g_j \quad z \in R_j$$

$R_1 \dots R_N$  is polyhedral partition of  $X$

Solution of any QP is piecewise affine in rhs of constraints

eg.  $l_1$  regularized problem

$$f(x) = \lambda \|x\|_1 \quad x^* \text{ is piece-wise affine in } \lambda$$

can build a look-up-table for  $\lambda, x^*$

$\Delta u_{\text{mpc}}$  (ie.  $K_j, g_j, R_j$ ) can be computed explicitly, off-line

## • MPC problem structure

MPC problem is high-structured

Hessian is block diagonal

equality constraint matrix is block banded

eg

$$L(X(t), U(t))$$

$X(t)$  is directly connected to

$$\begin{cases} U(t) \\ U(t-1) \\ X(t+1) \\ X(t-1) \end{cases}$$

$$X(t-1) \quad X(t) \quad X(t+1) \quad U(t) \quad U(t-1)$$

block of variables  $\Rightarrow$  banded

Use block-elimination to compute Newton step

## • Supply chain management

$n$  nodes (warehouses / buffers)

$m$  unidirectional links between nodes, external world

$X_i(t)$  is amount of commodity at node  $i$ , in period  $t$

$U_j(t)$  is amount of commodity transported along link  $j$

incoming / outgoing node incidence matrices

$$A_{ij}^{\text{in/out}} = \begin{cases} 1 & \text{link } j \text{ enters/exits node } i \\ 0 & \text{otherwise} \end{cases}$$

$$\text{dynamics} \quad X(t+1) = X(t) + A^{\text{in}} U(t) - A^{\text{out}} U(t)$$

buffer limits  $0 \leq X_i(t) \leq X_{\max}$  (could allow  $X_i(t) < 0$  to represent back order)

link capacity:  $0 \leq U_j(t) \leq U_{\max}$

$$A^{\text{out}}(t) \leq X(t) \quad (\text{can not ship out what's not on hand})$$

$$\text{objective: } \sum_{t=0}^{\infty} S(t) + w(X(t))$$

shipping cost + storage cost